#### COVER SHEET FOR TECHNICAL MEMORANDUM

of Documentation by Macro
Assemblers

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#### ABSTRACT

It is shown that macro assemblers can easily generate documentation, of a meaningful sort, which is, in at least one respect, superior to that which can be built into the 'self documenting', English-appearing languages such as ALGØL, NELIAC, etc.

Due to currently common limitations, however, this can be done only for very small programs. In particular, trivial difficulties with scan routines and unrealistic limitations on conditional assembly tend to contribute to the main problem of table overflow. It is found that an ALGØL-type table clean out feature would clear the way for practical use.

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of Documentation by Macro
Assemblers - Case 210

DATE: September 2, 1964

FROM: W. M. Keese, Jr.

## MEMORANDUM FOR FILE

#### Introduction

This note follows an interesting foray into automatic generation of meaningful documentation. The foray was undertaken in belief in the single document approach for a program, i.e., belief that a program listing, itself, should contain all of the documentation which the program requires\* - including even the raw content of its various levels of flow charts.

This means that the source deck of a large program should contain many levels of documentation. There should be considerable explanation of the segmentation. Much material about data structures, good paragraphization within segments, etc. There is much of interest in automatic aids for this higher area, but this exercise is not concerned with higher levels. Rather, it addresses itself to the documentation of individual source instructions, particularly those of a higher level than simple arithmetic.

It is often the case that a good percentage of the source instructions in any large program are involved with decision making, calling of subroutines, transfers of control, and the simple moving of data. Somewhere, to be sure, there are routines which actually do things to data - which use arithmetic, Boolean, shift, etc. instructions. But, the larger the program is, the smaller part of it that these instructions seem to comprise.

In fundamental level documentation, no programmer needs to be told what (e.g.) a compare instruction does, but he does want to know what it is that is being compared. Simple machine knowledge allows recognition of a decision, but only an intimate

<sup>\*</sup>This is not to say, however, that there should not be any other documentation.

knowledge of a program's data base may be able to provide knowledge of what decision is being made. Where subroutines are being called (or macros invoked), the situation is even worse, for no general knowledge is of any use in trying to follow the program. It can only be done if the call is accompanied by commentary that explains what is done.

This kind of commentary is vital to the feasibility of changing large programs and perhaps the most important factor in controlling the cost of such. Yet, experience tells us all that even the best of intentions are not always sufficient to assure that it is always correct and/or up-to-date.

When such documentation can be automatically generated, one is relieved of this large area of concern in programming control.

The enclosed example shows that it is easy, with a quite small macro package, to automatically generate good, meaningful commentary for the kinds of things outlined above -documentation about decisions and control, and explanation of the effect(s) of subroutines.

The possibility of this hinges upon the fact that all such actions can readily be performed by macros. These can, in fact, all be done with a standard macro package (see appendix). The use of such a package for decision and control is quite standard, and needs no explanation here. It is shown in section 1. that the calling of subroutines can and should be effected in a similar way, with no extra macro definitions on the part of the programmer. (An added advantage of treating all subroutines, in the source language of the calling program, as though they were macros is that proper modularity can be retained in the source program even when it must be violated in the actual machine program.)

An example of what can be done terminates this section. The actual macro package which creates it is included in the appendix.

Since this macro package is rather confusing when approached as a whole, an heuristic approach is used. Most of the main concepts arise in the development of the subroutine documentation macros, and this is done in some detail in sections 1. - 5. Sections 1 and 2 give the milieu from which the rest arcse. Section 3 shows the generation of simple comments, and variable substitution is added in section 4. In section 5,

expansion is made to occur word for word on a recursive basis. The other macros used in the sample are briefly noted in section 6. Section 7 notes limitations and draws conclusions.

Particular note should be taken of the difference between the grade of documentation produced by the macros appearing in 3, 4, and 5. In 3, one gets nothing but the expanded subroutine name. In 4, however, one can obtain a meaningful sentence, with operands plugged into their proper (English sentence structure) places. This is at once a major improvement in documentation over functional notation, such as is found in compilers. In order to expand these operand names, so that the reader may fully understand the documentation without continual reference to the noun list, one has to go to the greater effort shown in section 5.

The various examples happen to be done in IBMAP for the 7040-44. In this, "IRP" followed by an argument, begins an indefinite iteration on the subfields of that argument. Its scope is ended by the next "IRP". "IFF" and "IFT" are conditioned assembly controls. With the first, the following card is assembled only if false that the (S) symbol value of the two expressions in the variable field are equal. With "IFT" the following card is assembled only if true that they are equal. Matched parentheses, in a macro instruction, are stripped, but make the entire character sequence which they enclose into a single argument. The apostrophe is simply the concatenation character.

. •			•
	USES		PEND.CHAR.TD.ID.STRING)
	USES	GNC()(GET,NC)	
	USES		BLANK.CHARACTER)
	USES	SAVID()(SAVE,I	DENTIFIER, STATUS)
•	WITH V	ARIABLE DECLARAT	IONS OF THE FORM
	VAR	NC(NEXT, CHARAC	
<del></del>	VAR	PC (PEEK, CHARAC	TER)
•			
			E PROGRAM MAY THEN BE WRITTEN
•	IN ABB	REVIATED FORM, A	S BELOW, IN THE SOURCE DECK.
*			
#	LABEL	ANAME (APUSTRUP	HE-ENCLOSED, NAME)
#	GNC		
*	IF	NC.NEQ(=H00000	•)
#	THEN	A 3 4 4 5 5	
	GOTO	ANAME	
#	ELSE	NC	
•	APPND	NC	
*	GNC GOTO	NONSPC	
•		NUNSPL	
•	•••		
•	LABEL	BLNKF(BLANK,FO	ILIND)
	APPND	PC	·
•	LAC	HEAD.I	SET POINTER TO HEAD OF TREE.
	GNBC	TICAUT I	JET TO MICK TO MEND OF THEE
-	TRA	0.1	TRACE SPECIAL WORD TREE.
<del></del>	108		TANGE OF LOTAL HORD THELE
		· · · · · · · · · · · · · · · · · · ·	

<del></del> -				- 5 -	09
<del></del>					
		•		THIS WILL BE EXPA	NDED IN THE LISTING AS IS SHOWN HE
		•			
	00006		LABEL EQU •	ANAME (APOSTROPHE-	
			GNC	А	POSTROPHE-ENCLOSED NAME
					GET NEXT CHARACTER .
<del></del>				NC.NEQ(=H000000)	• IF NEXT CHARACTER NOT EQUAL TO =H00000*
01	<del></del>				
	<del></del>		THEN		• THEN
			GOTO	ANAME	
				AUDITE	GO TO APOSTROPHE-ENCLOSED NAME
			ELSE		• ELSE
			APPND	NC	APPEND NEXT CHARACTER TO ID
	<del></del>		GNC		STRING ,
			GOTO	NONSPC	GET NEXT CHARACTER ,
				HUNGFU	GO TO NONSPC ,
		<u> </u>			• • • • • •
02		•			
			LABEL	BUNKF (BLANK . FOUND	))
	00016	BLNKF	EQU ●	p	SLANK FOUND
			APPND	PC	APPEND PEEK CHARACTER TO ID
	<del></del>				STRING .
1	00003 10	0001	LAC GNBC	HEAD, I	SET POINTER TO HEAD OF TREE.
			GNDC		GET NON BLANK CHARACTER .
1	00000 10	0000	TRA	0,1	TRACE SPECIAL WORD TREE.
					<u></u>
		-			<u> </u>

## 1. Background

The automatic documentation features described below were not created in vacua. Rather, they were recognized as a free bonus coming with certain other features being built into a set of standards for assembly language program packages.

Two of these are relevant. First: each routine or segment of a large program should carry with it a list of first level subroutines used. Second: while the fact that a subroutine is called is a part of the logical structure of the calling program, the way in which it is called is part of the logical structure of the subroutine - not of the calling program; it is therefore better modularity that the source language of the calling program should specify merely the existence, not the method, of the calls.

This second consideration gives rise to the notion that a particular call should be effected in an assembly language merely by writing the name of the subroutine in the operator code field (with any variable list(s) in the variable field). Particular knowledge of how a given subroutine is called can then be localized to one macro definition. Indeed, since, even in real time programs, most subroutines can be called in some standard manner, a standard call defining macro could be used to define most of the subroutines' calls.

The inclusion of a list of first order subroutines used, however, relieves us of even this necessity. This list, itself, can be used to define all the standard calls. Only those special routines which efficiency or special difficulty demand to have uncommon calls need have their names defined as call creating macros.

## 2. The Elementry Uses Macro

We were first led, then, to the creation of a USES macro. Each routine or segment of a large routine contains in its head a set of USES cards. Each of these cards contains the word 'USES' in its operator field and the name of a subroutine in its variable field, along with some explanation of the routine's effect.

The original purpose of these cards was simply the display of program organization. An elementry extension creates standard call definitions.

The USES macro first asks whether the subroutine name has been previously treated. If it has, it does nothing, but if it has not, then it defines it to create a standard system call. At this stage of development, the macro reads (in 7040 IBMAP)

USES MACRØ NAME IFF NAME . = DEFIND USES1 NAME ENDM USES USES1 MACRØ NAME NAME'. DEFIND SET NAME MACRØ ARGLST, ERRT CALL NAME (ARGLST) ERRT ENDM NAME ENDM USES1

(It is intended that this package run under no created symbols.)

The creation a symbol from "NAME" by concatenation of a period is possible for us, since the standards in use at Bellcomm limit the length of a subroutine name to one character less than maximum symbol length. Even without this, one could still ask whether definition had occured by making the meaning of the created NAME macro depend on whether or not some symbol was set to USES or to NØRMAL. Under the USES mode, NAME (if defined) would set the remainder of USES to be null (save for terminal self restoration). The remainder of USES would elsewise create the standard definition. This (using an operator-synonymous operator) could be done with no expandingly greater use of macro skeleton space, but it would be somewhat slower.

## 3. Generation of a Primitive Comment

Since rudimentary documentation practices demand that each of the USES cards carry some explanation, it is but a simple step forward to automatically reproduce this comment whenever the subroutine is called. This will be shown here in a primitive manner, while the next section will go on to show an elementry plugging in of arguments.

One merely encloses the explanation in parentheses (to make it a single argument) and moves it over from the comment field to adjoin the name. USES is now made a 2 argument macro. If it finds that NAME has not been defined, then it passes both the name and the (reparenthesized) comment along to USES1 which now reads

USESI MACRØ NAME, CØMENT NAME' SET DEFIND NAME MACRØ ARGLST, ERRT CALL NAME (ARGLST)ERRT PMC ØN REM PMC ØFF ENDM NAME ENDM USESI	'CØMENT'
--	----------

In its definition of NAME, USES1 has added 3 card images. These turn on the printing of macro expansion cards, print the comment as a remark, and turn off the printing of macro generated cards: (as PMC is a control card, it, itself, will not be printed.)

In actual practice, NAME could be made shorter, saving skeleton space, by passing "CØMENT" along to another macro, which would do the three card expansion.

## 4. Comments With Substitutable Arguments

Whereas many subroutines have arguments, and their effect is on these arguments, a meaningful comment generator should plug the names of these arguments into the comments. This is a relatively easy addition.

In the following expansion, we start afresh and use slightly more levels, so as to conserve skeleton space. An argument list is added to the USES macro to make it possible to tell what words of the explanation should be substituted.

USES	MACRØ IFF USES1 ENDM	NAME(ALIST)CØMENT DEFIND=NAME'. NAME(ALIST)(CØMENT) USES	
USES1 NAME'. NAME	MACRØ SET MACRØ USES2 ENDM	NAME(ALIST)CØMENT DEFIND ARGLST,ERRT NAME(ARGLST)(ERRT) NAME	
NAME'.	MACRØ	ALIST, DUMMY1, DUMMY2 (CØMENT)	
USES2	MACRØ CALL NAME'. ENDM	NAME (ARGLST)ERRT NAME (ARGLST)ERRT ARGLST,ERRT USES2	
PRSEQ	MACRØ PMC REM PMC ENDM	TEXT ØN ØFF USES3	'TEXT'

If one were not sure of the ability to concatenate the period to NAME, the macro "NAME'." could be named by a created label, generated by turning them on, going an extra level in, and turning them off. Its name would then be passed on to USES2.

We have reached the level which is best explained by an example, save for a preliminary note. While USES2 and PRSEQ

were created mainly\* to conserve skeleton space (for each of the presumably many, NAME and NAME'. macros), the creation of the "NAME'." macro serves a very different function. Since "ALIST" is built into its formal parameter list, while the comment is built into its skeleton, any delimited phrases in "CØMENT" which are the same as delimited phrases in "ALIST" will be substituted at expansion time. When a call is made (NAME is invoked) the actual parameter list is sent along to "NAME'." and the appropriate substitution is made for PRSEQ. The two dummy arguments of "NAME'." are merely to avoid a falacious diagnostic regarding argument count which 7040 1BSYS likes to make.

Suppose now that ADDTØ has not previously been defined and that the following source card appears.

USES ADDTØ(A,B)(ADD 'A' TØ 'B')

This will expand into

IFF DEFIND = ADDTØ.

USES1 ADDTØ(A,B)(ADD 'A' TØ 'B')

ADDTØ. SET DEFIND

ADDTØ MACRØ ARGLST, ERRT

USES2 ADDTØ(ARGLST)(ERRT)

ENDM ADDTØ

ADDTØ. MACRO A.B.DUMMY1, DUMMY2

PRSEQ (ADD 'A' TØ 'B')

ENDM ADDTØ

Now consider a later point in the text where a card appears with "ADDTØ" in the operator field. It will expand as below.

ADDTØ (ALPHA, BETA) (source card)

USES2 (ALPHA BETA)()

CALL ADDTØ(ALPHA, BETA)

ADDTØ. ALPHA, BETA

PRSEQ (ADD ALPHA TØ BETA)

PMC ØN

REM ADD ALPHA TØ BETA

PMC ØFF

The desired expansion is achieved.

<sup>\*</sup>for some assemblers, PRSEQ is necessary. See section 5.

# 5. The Print Sequence Macro, PRSEQ

In section 4., we introduced the use of a macro PRSEQ, ostensibly to save space. This macro (although not as shown there) is actually necessary to some assemblers, and is desirable in order that another level of documentation be achieved.

In fact, the system shown in section 4, will not work, as shown, on the 7040-IBMAP version 3, due to a fault in the assembler's macro skeleton scan. This fault is asserted to be corrected in version 6, but we have not yet seen evidence of this. To wit, USES is supposed to build the card image

## PRSEQ (CØMENT)

into the skeleton of "NAME'.". Such a card is indeed generated while this skeleton is being built up, but the scan in question stops at the first blank, not keeping a parentheses count. Thus, while the card generated for the skeleton is

PRSEQ (ADD 'A' TØ 'B')

, the line which is actually entered into the skeleton is

PRSEQ (ADD

. One can readily imagine to what grevious errors the generation of such a line leads the assembler. Such a scan error is apt to be met in many assemblers, so a way around it is important.\*

Further, the kind of comment generated by the macros of section 4, is not all that one might hope for. To wit, it is not sufficient that the symbolic names of variables be plugged into the comments generated by invoking a subroutine; one would prefer to plug in an explanation of what those names stand for. Put more simply, it is desirable that the expansion of comments should be a recursive process - going down through each word of the expansion (expanding it in turn, etc.) until all elements have been expanded to some bottom level.

<sup>\*</sup>It actually suffices, for our assembler, to build the body of PRSEQ, as shown, into USES1's definition of NAME'.

This aim can be accomplished, and its accomplishment bypasses any trouble that assembler scan routines may have about blank characters. In the accomplishment, all definitions given in section 4, except that of PRSEQ and USES2, stand as given. The use of USES, on the other hand, changes. Specifically, one separates each of the words in the comment skeleton by commas.

Thus, instead of writing

USES ADDTØ(A,B)(ADD 'A' TØ 'B')

one writes

USES ADDT $\emptyset(A,B)(ADD,A,T\emptyset,B)$ 

The definition of PRSEQ, however, becomes far more complicated. It becomes necessary to add words, one at a time, to a line image. USES2 becomes

USES2 MACRØ NAME(ARGLST)ERRT
CALL NAME(ARGLST)ERRT
PRSET SENT.
NAME'. ARGLST,ERRT
PRLINE (,)
ENDM USES2

Of the two additions, "PRSET SENT." sets up printing routines to use a sentence type of format, while "PRLINE (,)" causes the printing of a final line image with a comma ending the line. The generated macro "NAME'." still invokes PRSEQ, sending the comment (with actual arguments substituted) as a single argument. (In our example, "ADDTØ." would generate "PRSEQ (ADD, ALPHA, TØ, BETA)".)

PRSEQ, itself, merely does an indefinite iteration on the sub-arguments of its one argument - sending each of them on to a high level concatenate macro, PRADD.

The high level concatenate macro merely tests each of its arguments received for further expandibility. If it is so expandable, then this is done; elsewise, it passes the argument on to an exact concatenate macro, PREADD.

This concatenate macro, PREADD, is the heart of the entire automatic documentation system. It is the accessible

member of a pair of macros invented by Ann L. Locicero [1], which makes it possible to build up a linear image from successive calls of a macro.

For purposes of explanation, let us imagine that the concatenate macro were spelled "CØNCA" while its partner were spelled "CØNCAT".

The original definitions would be

CØNCA	MACRØ CØNCAT ENDM	TEXT (TEXT) CØNCA
CØNCAT CØNCA	MACRØ MACRØ CØNCAT ENDM ENDM	TEXT MØRTXT (TEXT,MØRTXT) CØNCA CØNCAT

Each invocation of CØNCA invokes CØNCAT so as to build its existant argument into all future calls. For instance, imagine the above definitions to be given and then the source cards

CØNCA A CØNCA B CØNCA C

to appear. The complete expansion would be

CØNCA	CØNCA CØNCAT MACRØ CØNCAT ENDM CØNCA	A (A) MØRTXT (A,MØRTXT) CØNCA B	Orig. meaning of CØNCA established meaning 2 of CØNCA.
CØNCA	CØNCAT MACRØ CØNCAT ENDM CØNCA CØNCA	(A,B) MØRTXT (A,B,MØRTXT) CØNCA C (A,B,C)	Meaning 2 of CØNCA established meaning 3 of CØNCA.

<sup>[1]</sup> Locicero, A. L.: "Linear Accumulation in Sequential Macros"; unpublished, Bellcomm, Inc.

CØNCA MACRØ MØRTXT
CØNCAT (A,B,C,MØRTXT)
ENDM CØNCA

Meaning 3 of CØNCA established meaning four of CONCA.

In actual practice, the partner macro also does bookkeeping to make sure that maximum legal argument length is not exceeded. This tends to make for an excessive amount of complication and could presumably be avoided in more adequate macro assemblers.

The PRSET routine uses its parameter to chose one of a set of actual print out macros used by PRLINE. It essentially sets a switch through the use of the operator-synonym operator.

The print line macro, PRLINE, sends its argument (if any) along to PREADD (properly parenthesized so that a separator comes out as a substitutable argument at the bottom level). It then invokes a print current line image macro, PR5., to cause the actual print. PR5. is also used by the concatenate partner macro, PR8., in case of line overflow.

The print current line image macro, PR5., changes the meaning of the concatenate partner, PR8., to that of a prepare for actual print macro, PR6., invokes the concatenate macro (which now invokes PR6. with all of its previous arguments built into its parameter), and then restores the concatenate macro and its partner to their initial meanings. The prepare for actual print macro, PR6., receives the built up line image as a single parameter (held together by its parentheses) from the prevertedly used concatenate macro, and sends it along through the dummy macro (which has been set by PRSET to be synonymous with one of the actual print macros). As PR6. has not parenthesized its single argument, this now comes apart into its several components which can be spaced apart from each other by the actual print macro (which always devolves into turning on macro print out, generating some remark card, and turning off macro printout).

Some complication could be done away with here if only it were possible to inbed blanks in the remark skeleton in the first place. Those having assemblers with such capability will find it easy to simplify the example.

# 6. Conditionals, Go To Statements, and Other Declarations

ALGØL-like conditional statements are easily put into an assembler in any number of ways. The if-then-else package used in the preceding example is an off-the-cuff version and could stand considerable improvement, particularly in the recognition of single word alternatives.

As a generator of documentation, it is an easy matter to have the various relational macros directly generate a remark line, plugging in variable names. By making use of the print package, it is possible to expand the names of the variables, exactly as is done in generating subroutine comments.

The assembly operations of the package are perfectly straightforward. The IF macro merely turns on created symbols and invokes an appropriate relation macro, such as GTR. This turns off created symbols, prints a comment, creates an appropriate comparison triplet, and sets up meanings for THEN & ELSE.

The design objective was that the third member of the comparison triplet never have to contain a comparison transfer. Any assembly reached by the other alternative is to be remote. The particular system uses the location counter choosing facilities available in IBSYS, but any remote facility could serve as well. The use of the macros USEN and USEP is purely for conformity with other standards in use where this package was written. (These two macros make up for the fact that USE X and USE PREVIØUS fail to operate in a push down fashion.) If-Then-Else allows specifying either the arithmetic or logical accumulator but does not presently guard against trying to load one from the other.

The GØTØ macro is thrown in merely to allow automatic commentary. AS LABEL declarations (unlike USES & VAR) need not precede the use of the associated identifiers, there is no guarantee that GØTØ will expand the name of the destination. (It will expand it if and only if it is a transfer backwards.) "GØTØ" is used rather than redefining "TRA" because the present automatic documentation facilities cannot handle indexed addresses.

LABEL and VAR are essentially dimensioning operators. The first defines a symbol and an associated comment, and

produces that comment as a paragraph heading, followed by a colon (".."). VAR remotely assigns a full word variable and associates a commentary name with its symbolic name.

If a symbol which has been defined by either of these appears as a parameter, then its expanded name will appear in the generated commentary.

## 7. Conclusions

The example shown in the introduction shows that a high grade of automatic documentation is possible. However, it is not practical at this level, for one runs into limitations on table sizes.

Programs using this full level of expanded documentation run into overflows at between two and fifteen pages of expansion, depending on content.

Two possible courses of action are open: one can settle for the level of documentation generated in section 4., or one can move to repair the assembler deficiencies which cause overflow.

The section 4. level of documentation is marginally adequate. If name lists are particularly small or orderly, or if documentation requirements elsewise do not require explanation of all operands with each use, then the documentation created at this level is entirely satisfactory. The generation is extremely simple and cheap, and gives no trouble with table sizes. At almost no cost, and at some lessening of normal programmer effort, one can assure that documentation is always accurate and current. Others will find that this grade of documentation is just not good enough.

As to repairing the assembler characteristics which make the full course impractical, there are three things which would help a bit, and one which would entirely suffice:

Proper scan facilities in the building of macro skeletons would allow the expansion of comments to be done in many less steps - resulting in considerable saving in the number of skeletons generated.

Proper handling of end of line overflow on generated remark cards would relieve the macro package of considerable bookkeeping bother.

Greater table sizes would help a bit.

This last measure, however, would only be a small stop-gap. What is needed is some way to clear the tables of entries no longer needed.

While one could evolve all sorts of complicated schemes for this, the ideal solution seems to be the adoption of an ALGØL - like block structure - at least for macro definitions. In this structure, one could define BEGIN-END pairs (which could be nested). Every time a BEGIN was encountered, the condition of the tables could be marked. On meeting the corresponding END, the condition could be restored. This can be implemented merely by the addition of two words to each macro skeleton (with the addition of some elsewise vacuous skeletons for operator-synonym operations).

1031-WMK-mat

Attachment Appendix W. M. Keese, Jr.

# BELLCOMM, INC.

# APPENDIX

The following pages include the macro package used in the generation of the example shown in the introduction.

TEXT	WRITE			
·	•		AUTOMATIC	DOCUMENTATION WM KEESE 8/15/64
	. •			USEN-USE
		· · · · · · · · · · · · · · · · · · ·		· · · · · · · · · · · · · · · · · · ·
	FO.	MACRO		INITIAL RESTORATION MACRO
	<u>F1.</u>	DPSYN	FO.	RESTORE SELF TO INITIAL.
	USEP	OPSYN . USE	F.USE	SET 'USEP' TO GIVE FAULT, USE BLANK LOCATION COUNTER.
	·	ENDM	FO.	OJE JUNIO EGONTEN
	F1.	OPSYN	F0.	
	F2.	MACRO	X	USED TO HOLD REDEFINITIONS OF F1. TO 1
·	F1.	OPSYN	X	RESTORE TRACE FORWARD MACRD.
	USEP	DPSYN	X.	SET 'USEP' TO CURRENT USE NAME,
	····	USE	_X	USE CURRENT USE NAME LOCATION CNTR.
	<u></u>	ENDM,	F2.	
	F.USE	MACRO		ILLEGAL 'USEP' FAULT
	F.U3E	PMC	ON	ILLEUAL USEF FAULTA.
		SPACE	1	
		REM	*******	
		REM	*ATTEMPT	TO USE -USEP- WITH NO -USEN- IN EFFECT- IGNORE
	-	REM	****	
		SPACE	1	'
		PMC	OFF	
	•	ENDM	F.USE	
	USEN.	MACRO	CL	WORKING USE NEW MACRD
	CL	OPSYN	F1.	SET CREATED LABEL FOR PERM TRACE BA
	USEP	OPSYN	F1.	SET 'USEP' FOR CURRENT TRACE BACK,
		_, _,,,,		UPDATE TRACE BACK MACRO
	F1.	MACRO	· · · · · · · · · · · · · · · · · · ·	BUILD CL NAME INTO MACRO TABLE
		F2.	CL	TO REDEF 'USEP' AND RESTORE LOC CTR
		ENDM	F1.	(SEE F2. FOR RESTORATION.)
	•			
		USE	CL	USE NEW LOC CTR.
	<del></del>	ENDM	USEN.	
	•			
	USEN	MACRO	ANYLAB	"USE NEW! MACRO
	•	ORGCRS		UNLESS A NAME IS GIVEN, CREATE DNE.
		USEN.	ANYLAB	DO THE WORK,
		NOCRS		TURN OFF SYMBOL CREATOR.
	•		HCEM	
		ENDM	USEN	
	• USED	ENDM		HICCOL THITTALLY BARRIPES CAME
	USEP		F.USE	'USEP' INITIALLY PRODUCES FAULT.
		ENDM		'USEP' INITIALLY PRODUCES FAULT.
		ENDM		'USEP' INITIALLY PRODUCES FAULT.
		ENDM		'USEP' INITIALLY PRODUCES FAULT.

:				IF-THEN-ELS
	* AAC	DPSYN	CLA	LOAD ARITHMETIC AC.
······································		UF J III	· · · · · · · · · · · · · · · · · · ·	LUAU ARTITMETTO Acce
	DEFOR	MACRO	RM-CL-CLD	BUILD LABELS + REMOTE TYPE INTO -OP
	OP.	MACRO	T	MAKE -DP.(TYPE)- MEAN
		T. RM	•	INVOKE TYPE OF REMOTETYPE.
		ENDM	OP.	
		ENDM	DEFOP	
	• ,			
	ELSE_	MACRO		ELSE
		PMC	ON	PRINT REMARK,
		REM		• ELSE
		PMC	OFF	
	<del></del>	OP.	<u> </u>	DO OP.(ELSE).
		ENDM	ELSE	
				POWE TO 1
	EQU.	MACRO	A,B,CLT,CLD	EQUAL TD
		NOCRS	71.10 175 11-101	WASHIN MPHARY
		PRSTAT	FLAG. (IF, A(=)B)	
		REL		BUILD LABELS AND COMPARE TRIPLET.  BUILD LABELS AND 'REMOTE THEN' IN O
		DEFOP	RT,CLT,CLD	BOILD FADERS WAS TREATE THEM. THE ME
*····	•	ENDM		
	E RE	MACRO	CLF.CLD	ELSE FOR REMOTE ELSE
	E	SET	1	MARK 'ELSE' DONE,
	•••	USEN	L	START NEW LOCATION COUNTER,
	CLF	EQU	•	DEFINE FALSE LABEL.
	VE-	ENDM	F.RF	VET SITE TIMES BROKET
	E-RT	MACRO	CLT.CLD	ELSE FOR REMOTE THEN
		TRA	CLD	GENERATE TRANSFER TO END OF COND.
		USEP		USE PREVIOUS LOCATION COUNTER.
		ENDM	E.RT	
	GEQ.	MACRO	A,B,CLT,CLD	GREATER THAN OR EQUAL TO
		NOCRS		
		PRSTAT		ATER, THAN, DR(=)B) PRINT REMARK,
		REL		T.A.B CREATE LOAD AND COMPARE TRIPLET.
		DEFOP	RT,CLT,CLD	BUILD LABELS AND REMOTE THEN IN OP.
····		ENDM	GEQ.	
	• .			
	GTR.	MACRO	A.B.CLT.CLD	GREATER THAN.
		NOCRS	45 A ORE	APPR TILL RE BREWE BEMARK
		PRSTAT		ATER, THAN, B) PRINT REMARK.
		REL	RT-CLT-CLD	1,A,B CREATE LOAD AND COMPARE TRIPLET,
		DEFOR		RUILD LARELS AND REMOTE THEN INTO O
		ENDM	GTR.	

	•			IF-THEN-ELSE
<del></del>	•			
	1F	MACRO	A.REL.B	lf
	••E	SET	0	MARK 'ELSE' UNDONE,
		ORGCRS	<u>·</u>	TURN ON CREATED SYMBOLS TO CREATE
		REL'.	A,B	2 LABELS AND DO APPROPRIATE RELATION,
	-	NOCRS		MAKE SURE CREATED SYMBOLS OFF.
-		ENDM	1F	
	•			
•	LAC	OPSYN	CAL	LOAD LOGATCAL AC
<del></del>	•			
	LEQ.	MACRO	A,B,CLE,CLD	LESS THAN OR EQUAL TO
		NOCRS		
	•	PRSTAT	FLAG. (IF, A, LESS	THAN, OR, (=) TO, B) PRINT REMARK,
	·	REL		.A.B CREATE LOAD AND COMPARE TRIPLET.
		DEFOP	RE,CLE,CLD	BUILD LABELS AND REMOTE ELSE INTO OP.
<u> </u>	·	ENDM	LEQ.	
	•	*****		
	LOAD	MACRO	REG.X	LOAD REGISTER
	•	IFF	REG=X	IF OPERAND /= REGISTER
	<del></del>	REG	<u> </u>	THEN DO APPROPRIATE LOAD.
		ENDM	LOAD	
	LSS.	MACRO	A.B.CLE.CLD	LESS THAN
	F22•	NOCRS	Aphyllegulu	FE22 I LINIO ●
······································	· · · · · · · · · · · · · · · · · · ·	PRSTAT	ELAC-ITE-A-LESS	THAN, B) PRINT REMARK,
•		REL	AAC_CAS_CIF.CLF	.A.B CREATE LOAD AND COMPARE TRIPLET.
		DEFOP	RE, CLE, CLD	BUILD LABELS AND REMOTE ELSE INTO DP.
		ENDM	LSS.	DOLLO ENDESS MIS MANAGEMENT
	+	<u> </u>		
	NEQ.	MACRO	A.B.CLE.CLD	NOT EQUAL TO
		NOCRS		
		PRSTAT	FLAG. (IF, A, NOT,	EQUAL, TO, B) PRINT REMARK,
		REL	LAC, LAS, ++2, CLF	A,B CREATE LOAD AND COMPARE TRIPLET,
		DEFOR	RE.CLE.CLD	
		ENDM	NEQ.	
	<del></del>			
•				
				•
** · · · · · · · · · · · · · · · · · ·				

	•			
· .				CREATE LOAD AND COMPARE TRIPLET
	REL	MACRO	REG. COMP. LG1	
		LOAD	REG.A.	
		COMP	В	SET UP COMPARISON,
		TRA	LGTR	IF GREATER, GO TO GREATER LABEL.
		TRA	LEQ	IF EQUAL, GO TO EQUAL LABEL.
		ENDM	REL	ASSUME LESS THAN CASE FOLLOWS.
	•			•
	THEN			THEN
		PMC	ON	PRINT REMARK.
		REM		• THEN
		PMC	OFF	
		OP.	<u> </u>	DO OP. (THEN).
		ENDM	THEN	
<del></del>				
	T.RE	OPSYN	NULL	THEN FOR REMOTE ELSE (= NULL)
· · · · · · · · · · · · · · · · · · ·	T.RT	MACRO	CLT, CLD	THEN FOR REMOTE THEN
	1000	USEN	0217020	USE REMOTE LOCATION COUNTER.
•	CLT	EQU	•	DEFINE TRUE LABEL.
	CLI	ENDM	TRT	DEI INC TRUE ENDEE
	•			
	RE	MACRO_	CLE.CLD	END OF CONDITIONAL FOR REMOTE ELSE
		IFT	••E=0	IF THERE WAS NO 'ELSE'
	CLE	EQU	_CLD	DEFINE ELSE-LOC = DONE-LOC.
		IFF	••E=0	IF THERE WAS AN 'ELSE'
		TRA	CID	TRNSFER OUT OF THE REMOTE.
		IFF	••E=0	
		USEP	••6-0	AND END THE REMOTE SECTION.
	CLD	EQU	•	DEFINE END OF COND. LABEL.
	CEU	ENDM	RE	DELINE FUD OF COURS TABLES
	•	ENUM		
	RI	MACRO	CLT.CLD	END OF CONDITIONAL FOR REMOTE THEN
	CLD			DEFINE END OF COND. LABEL.
	CLD	EQU	• RT	ARITHE CHA OL COMPS ENDERS
		ENDM		
	_	MACRO	•	END OF CONDITIONAL.
			ON	PRINT REMARK,
		PMC REM	ON	FRINI REMARKS
			OFF	
		PMC OP	UFF	DO DP. (END OF CONDITIONAL).
				UU UF LUIU UF LUIUL LUIRLA
		ENDM	•••	

			•	
 -PRINT	MACROS		•	PR
•				<u></u>
 <b></b>		THESE MACROS A	RE USED TO BUILD UP AND PRINT LINE	
•			OMATIC GENERATION OF COMMENTS.	
 •		THE MAIN 'ENTR	IES ARE	
		,,,,,,,,,,,,,,,,,,,,,,,,,,,,,,,,,,,,,,,		
 • .	PRADD		APPEND PHRASE TO LINE	
<u> </u>	PREADD		APPEND EXACT PHRASE TO LINE	
• ,	PRLINE	(ANY PUNCTUATION)		
 	PRSEQ	(PHRASE SEQUENCE)		
 •	PRSTAT	TYPE (MESAGE) PUNCT		
	PRSET	TYPE	SET LINE TYPE .	
•	1 1			
 		WHERE 'TYPE' STAN	IDS FOR EITHER 'HEAD.'. 'FLAG.'.	
●,		OR 'SENT.'.		
 			·	
 •		THE OTHER INTE	RNALLY USED PARTS ARE	
	<u>.</u>			
•	PR1.		ACTUAL HEADER PRINT	
 	PR2.		ACTUAL FLAGGED PRINT	
 . • .	PR3.		ACTUAL SENTENCE PRINT	
•	PR4.		ACTUAL PRINT DUMMY	
•	PR5.		PRINT CURRENT LINE	
•	PR6.		STRIP PARENS FOR ACTUAL PRINT	
 •	PRT.		NORMAL APPEND PARTNER	
•	PR8.		APPEND PARTNER	
•	PR11.		SET HEADER PRINTING	
 	PR12.		SET FLAGGED PRINTING	
•	PR13.	7	SET SENTENCE PRINTING	
		:	* * * * * * * * * * * * * * * * * * *	
•				
 PR1.	MACRO	A.B.C.D.E.F.G.H	CTUAL HEADER PRINT	
	PMC	ON		
 •	REM	' .	A B C O D E FIG	
	PMC	OFF		
	ENDM	PR1.		
• .		<del></del>		
 PR2.	MACRO	A.B.C.D.E.F.G.H	CTUAL FLAGGED PRINT	
	PMC	ON		
	REM		• 141 181 1C1 1D1 1E1 1F11G1	
•	PMC	OFF		
 	PR13.		RESET PRINT OUT FOR NORMAL LINES	
•	ENDM	PR2.		
PR3.	MACRO	A,B,C,D,E,F,G	ACTUAL SENTENCE PRINT	
	PMC	ON		
	REM		141 181 1C1 1D1 1E11F1	
 	PMC	OFF		
,	ENDM	PR3.		
 •	-			
PR4.	OPSYN	PR2.	ACTUAL PRINT DUMMY	
•			(SWITCHED BY PRSET)	

	PR5.	MACRO		PRINT CURRENT LINE.
	PR8.	DPSYN	PR6.	CHANGE APPEND PARTNER TO MEAN STRIP PARENS FOR ACTUAL PRINT
		PREADO		AND INVOKE IT.
	PRWCT.	SET		RESET WORD COUNT.
,	PR8.	DPSYN	PR7.	RESET APPEND PARTNER.
	PREADD		PRO.	RESET APPEND EXACT WORD.
		ENDM	PR5.	
	PR6.	MACRO	TEXT	STRIP PARENS FOR ACTUAL PRINT
		PR4.	TEXT	SEND SEPARATED ARGS THROUGH
	_	ENDM	PR6.	ACTUAL PRINT DUMMY.
	PR7.	MACRO	TEXT	NORMAL APPEND PARTNER
	PREADD	MACRO	MORTXT	REDEFINE APPEND EXACT WORD
		PR8.	(TEXT, MORTXT) PREADD	TO INVOKE PARTNER WITH TEXT BUILT IN
	PRWCT.	SET	PRWCT.+1	UP WORD COUNT.
		IFT		IF WORD COUNT = WORD LIMIT
		PR5.	PR7	THEN PRINT CURRENT LINE.
	•			
	PR8	OPSYN	PR7	APPEND PARTNER (SET TO NORMAL)
	•			
	PR11.		003	SET FOR HEADER LINES MAKE-ACT. PRINTACT. HEAD. PRINT-,
	PR4.	DPSYN	PR1.	A = WORD LIMIT.
		ENDM	PR11.	D. S. RURY LIGHT
	PR12.	MACRO		SET FOR FLAGGED PRINT
	PR4.		PR2	MAKE ACTUAL PRINT = FLAGGED PRINT.
	PRWLM.		6	6 =. WORD LIMIT.
		FNDM	PR12	
	•			
	PR13.	MACRO		SET FOR SENTENCE LINES
	PR4.	DPSYN	PR3.	MAKE-ACT. PRINTACT. SENTENCE PRINT-
	PRWLM.		5	5 = WORD LIMIT.
	· .	ENDM :	PR13.	
	PR14.	MACRO OP	OP	
		ENDM	PR14.	

	•			
	•			
	PRADD	MACRO	TEXT	ADD PHRASE TO LINE
		IFT	DEFIND=TEXT.	IF TEXT IS DEFINED
		PR14.	TEXT.	THEN EXPAND IT
		IFF	DEFIND=TEXT.	ELSE
		PREADD	(TEXT)	ADD IT DIRECTLY TO LINE.
		ENDM	PRADD	,
	• ;			
	PREADD	OPSYN	PR8.	APPEND EXACT PHRASE TO LINE
	•			
	PRLINE	MACRO	PUNCT	PRINT LINE
	PRWLM.	SET	PRWLM.+1	UP WORD LIMIT FOR PUNCTUATION.
		PREADD	((PUNCT))	ADD ANY PUNCTUATION,
	PRWLM.		PRWLM1	RESTORE WORD LIMIT.
		PR5.		PRINT CURRENT LINE IMAGE.
		ENDM	PRLINE	
	•			
	PRSET	MACRO	TYPE	SET LINE TYPE
		IFT	TYPE=HEAD.	(DECDDES -TYPE- AND
		PR11.	TITE TIERDS	CALLS APPROPRIATE INITIALIZER.
		IFT	TYPE=FLAG.	The state of the s
		PR12.	THE TEAGS	
		IFT	TYPE=SENT.	
		PR13.	, , , , , , , , , , , , , , , , , , , ,	
		ENDM	PRSET	
	•	LITOIT	, Koe i	
······································	PRSEQ	MACRO	COMENT	APPENT PHRASE SEQUENCE
	. 1/364	IRP	COMENT	Mile Citi I illinam omdominamo
		PRADD	(COMENT)	APPEND PHRASE TO LINE.
	r,	IRP	1001161117	THE RESTRICT OF THE PARTY OF TH
<del></del>		ENDM	PRSEQ	
	•	CHON	. 11464	
<del></del>	PRSTAT		TYPE, MESAGE, PU	NCT
	INJIMI	PRSET	TYPE	
<del></del>		PRSEQ	(MESAGE)	
	*	PRLINE	(PUNCT)	
	<del></del>	ENDM	PRSTAT	
		CHUM	FRSIMI	

		,		us
	GOTO	MACRO TRA	LABEL LABEL	
		PRSET	SENT.	
		PREADD	<del></del>	
		PREADD PRADD	TO LABEL	
		PRLINE	(,)	,
		ENDM	GOTO	
	• .	CHUM		*
	LABEL	MACRO	NAME COMENT	
		PMC	ON	
	NAME	EQU	•	
		PMC	OFF	
	NAME .	SET	DEFIND	MARK NAME AS BEING DEFINED.
	NAME".	MACRO		DEFINE ASSOCIATED EXPANSION MACRO
		PRSEQ	(COMENT)	PRINT OUT PHRASES IN MACRO.
		ENDM	NAME .	
	`.	PRSET	HEAD.	SET PRINTER TO PRINT HEADING.
		NAME'.		PUT COMMENT PHRASES IN LINE IMAGE,
<del></del>		PRLINE	()	PRINT LINE FOLLOWED BY COLON.
		ENDM	LABEL	
<del></del>	*****	***	CO / AL TOTA CONTAIT	
	USES	MACRO	SR(ALIST)COMENT SR'.=DEFIND	IF SUBROUTINE MACRO NOT DEFINED
<del></del> -		USES1.	SR(ALIST)(COMENT)	THEN DEFINE IT
		ENDM	USES	THE OUT THE IT
	•			
	USES1.		SRIALISTICOMENT.	NOTE THAT SR IS DEFINED.
	SR'.	SET	DEFIND	DEFINE SR NAME
	SR	MACRO USES2.	ARGLST, ERRT SR(ARGLST) (ERRT)	TO INVOKE STANDARD CALL.
		ENDM.	SR	TO THYUNE STANDARD DALL.
	SR*.	MACRO	ALIST, DUMY1, DUMY2	DEFINE AUGUMENTED NAME
	3K • •	PRSEO	(COMENT)	TO APPEND ALL COMMENT WORDS.
		ENDM	SR'.	
		ENDM	USES1.	
	•		:	
	USES2.	MACRO	SR(ARGLST) ERRT	*
		CALL	SR(ARGLST)ERRT	PRODUCE DESIRED STD. CALL,
	·	PRSET	SENT	READY PRINT MACROS FOR SENTENCE.
		SR'.	ARGLST, ERRT	PUT COMMENT WORDS IN LINE IMAGE.
	· · · · · · · · · · · · · · · · · · ·	PRLINE		PRINT LINE WITH COMMA.
	v 7 .	ENDM	USES2.	•
	VAR	MACRO	NAME, COMENT	
	7711	USEN	STOR	
	NAME			
		USEP		
	NAME'.		DEFIND	
	NAME .	MALKU		
		PRSEQ	(COMENT)	
			(COMENT)	

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<u></u>			· · · · · · · · · · · · · · · · · · ·				
	•			· · · · · · · · · · · · · · · · · · ·			
	•					5	
	DEFSYM	4 MACRO	LIST	· /			
	LIST	IRP SET	LIST				
<del></del>	DEF.	SET	DEF.+1				
	<i>UL</i> : •	IRP	DEFAVE				
		ENDM	DEFSYM				
37200	DEF.		16000	· · · · · · · · · · · · · · · · · · ·		·	
	•				•		
<del></del>		DEFSYM	(AAC.DEFI	ND.LAC)			<del></del>
		DEFSYM	(HEAD. + FL	.AG., SENT.)			
<del></del>			•				
00001	ī	SET	1	. •			
<u> </u>	•						
	0000 APPND						
00000 10		<del></del>					
	OOOO GNBC	•					
00000 LU	OOOO HEAD	•	γ		į		
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•	ONE HEADS EACH ROUTINE OR SEGMENT WITH A "USES" LIST.E.G.					
•		400404644644				
<del></del>	USES		PPEND, CHAR, TO, ID, STRING)			
	USES	GNC()(GET,NC)				
	USES		N.BLANK, CHARACTER)			
•	USES	SAAID()(SAAE*	IDENTIFIER, STATUS)			
•	WITH VA	ARIABLE DECLARA	TIONS OF THE FORM			
	VAR	NC (NEXT, CHARA				
	VAR	PC (PEEK, CHARA	CTER			
•		THE TEXT OF T	UE DOCCAM MAY TUEN DE UNITTEN			
<del></del>	TN ARRI		HE PROGRAM MAY THEN BE WRITTEN AS BELOW, IN THE SOURCE DECK.			
	TIA WOOL	CATMICD LOWN	AS DEEDWY IN THE SUUNCE DEERS			
•	LABEL GNC	ANAME (APOSTRO	PHE-ENCLOSED, NAME)			
•	1F	NC , NEQ ( = H0000	0*)			
	THEN GOTO	ANAME				
•	ELSE	MINATE				
	APPND	NC				
	- GNC					
•	GOTO	NONSPC				
	***					
•	LABEL	BLNKF (BLANK, F	OUND)			
•	APPND	PC				
	LAC	HEAD.1	SET POINTER TO HEAD OF TREE.			
•	GNBC		·			
	TRA		TRACE SPECIAL WORD TREE.			
·						
		:				
<del></del>	<del> </del>					
····						

					09/04/
		•		THIS WILL BE EXPA	ANDED IN THE LISTING AS IS SHOWN HERE.
00006		-	LABEL	ANAME (APOSTROPHE-	ENCLOSED, NAME)
		ANAME	EQU •	A	APOSTROPHE-ENCLOSED NAME
				NC,NEQ(=H00000°)	GET NEXT CHARACTER .
		<del></del>	<del></del>	·	• IF NEXT CHARACTER NOT EQUAL TO =H00000*
		<del></del>	THEN		• THEN
			GOTO	ANAME	GO TO APOSTROPHE-ENCLOSED NAME ,
			ELSE		• ELSE
		<del></del>	APPND	NC	APPEND NEXT CHARACTER TO ID
			GNC		GET NEXT CHARACTER .
			GDTO	NONSPC	GO TO NONSPC .
	-		• • •		*
		•	LABEL	BLNKF(BLANK, FOUND	ni
00016		BLNKF	EQU •		BLANK FOUND
			APPND	PC	APPEND PEEK CHARACTER TO ID
00003	10001		LAC GNBC	HEAD, I	SET POINTER TO HEAD OF TREE,
00000	10000		TRA	0.1	GET NON BLANK CHARACTER , TRACE SPECIAL WORD TREE.
				-	
***					
				·	